Pseudo standard entanglement structure cannot be distinguished from standard entanglement structure

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Beyond IID 2022 arXiv:2203.07968

# **Brief Summary**

- Entanglement structures (ES)
  - is a possible structure of quantum composite system in General Probabilistic Theories (GPTs)
  - is not uniquely determined even if local structures are equivalent to the standard quantum theory.
- Pseudo standard entanglement structure (PSES) :
  - available for small error verification of all maximally entangled states
  - self-dual (= saturated situation of projectivity)
- Problem : Is there any possibility of PSES except for the standard entanglement structure?
  - infinitely many!
  - even if the PSES attains perfect discrimination of non-orthogonal states
- Variety of ES with group symmetry
  - global unitary symmetry determines ES as the standard entanglement structure (SES)



## 1 Preliminary:Brief Introduction to GPTs and Entanglement Structures

**2** Our Setting and Results

**3** Summary and open problems

## Outline

## 1 Preliminary:Brief Introduction to GPTs and Entanglement Structures

Our Setting and Results

**3** Summary and open problems

# Preliminary: Definition of Quantum Theory

Model of standard quantum theory on composite system  $\ensuremath{\mathcal{H}}$ 

- State  $\rho$  (density matrix)
  - $\rho \in \mathcal{T}_+(\mathcal{H})$  $\operatorname{Tr} \rho = 1$
- Measurement  $\{M_i\}_i$  (POVM)
  - $\stackrel{\bullet}{\Rightarrow} \frac{M_i \in \mathcal{T}_+(\mathcal{H})}{\operatorname{Tr} \rho M_i \ge 0} \ (\forall \rho \in \mathcal{T}_+(\mathcal{H}))$ 
    - $\sum_i M_i = I$

Assumption and Notation  $\dim(\mathcal{H}) = d$   $\mathcal{T}(\mathcal{H}) : \text{set of Hermitian}$ matrices on  $\mathcal{H}$   $\mathcal{T}_{+}(\mathcal{H}) : \text{set of positive}$ semi-definite matrices on  $\mathcal{H}$ 

- Probability to get a measurement outcome
  - outcome i is measured with probability  $\operatorname{Tr} \rho M_i$

 $\rightarrow$  the model of quantum theory is defined by  $\mathcal{T}_+(\mathcal{H})$ 

## Preliminary: Definition of the Model of Quantum Theory

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  - $M_i \in \mathcal{T}_+(\mathcal{H})$  $\Leftrightarrow \operatorname{Tr} \rho M_i \ge 0 \ (\forall \rho \in \mathcal{T}_+(\mathcal{H}))$  $\succ \sum_i M_i = I$

Essential Points  
• 
$$\operatorname{Tr} \rho M_i \ge 0$$
  
•  $\sum_i \operatorname{Tr} \rho M_i = 1$   
 $\rightarrow \{\operatorname{Tr} \rho M_i\}_i \text{ is probability distribution}$ 

- Probability to get a measurement outcome
  - outcome i is measured with probability  $\operatorname{Tr} \rho M_i$

 $\rightarrow$  the model of quantum theory is defined by  $\mathcal{T}_+(\mathcal{H})$ 

## Preliminary: Definition of the Model of GPTs

Model of GPT with positive cone  $\mathcal{K}(\subset \mathcal{T}(\mathcal{H}))$ 

- State  $\rho$ 
  - $\rho \in \mathcal{K}$  $\operatorname{Tr} \rho = 1$
- Measurement  $\{M_i\}_i$ 
  - $M_i \in \mathcal{K}^*$  $\Leftrightarrow \operatorname{Tr} \rho M_i \ge 0 \ (\forall \rho \in \mathcal{K})$  $\succ \sum_i M_i = I$

	Essential Points	
	• $\operatorname{Tr} \rho M_i \ge 0$	
Y	• $\sum_{i} \operatorname{Tr} \rho M_i = 1$	
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	probability distribution	J
	\ \	

- Probability to get a measurement outcome
  - outcome i is measured with probability  $\operatorname{Tr} \rho M_i$
- $\rightarrow\,$  a model of GPTs is defined by  ${\cal K}$

Remark Actually, GPT deals with more general models

# Entanglement Structure

Entanglement Structure (ES)

• ES is a possible structure of the model of the (bipartite) composite system of quantum subsystems in GPTs

Notations

•  $\mathcal{T}_{+}(\mathcal{H}_{A}) \otimes \mathcal{T}_{+}(\mathcal{H}_{B}) := \{\sum_{i} \rho_{i}^{A} \otimes \rho_{i}^{B} \mid \rho_{i}^{A} \in \mathcal{T}_{+}(\mathcal{H}), \rho_{i}^{B} \in \mathcal{T}_{+}(\mathcal{H})\}$ •  $\mathcal{K}^{*} := \{y \in \mathcal{T}(\mathcal{H}) \mid \operatorname{Tr} xy \geq 0 \ \forall x \in \mathcal{K}\}$ 

- ES is not uniquely determined to the Standard Entanglement Structure (SES) corresponding to  $\mathcal{T}_+(\mathcal{H}_A\otimes\mathcal{H}_B)$
- $\mathcal{T}_+(\mathcal{H}_A)\otimes\mathcal{T}_+(\mathcal{H}_B)\subset\mathcal{K}\subset(\mathcal{T}_+(\mathcal{H}_A)\otimes\mathcal{T}_+(\mathcal{H}_B))^*$ 
  - Local Tomography
  - No-Signaling Principle
  - Availability of separable states and separable effects
- it is important problem to derive the SES from ESs
- This talk shows the existence of "pseudo structures" of the SES.

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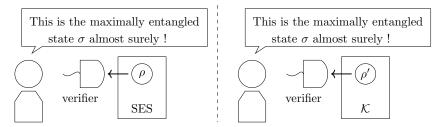
**3** Summary and open problems

## Pseudo Structures

Pseudo Standard Entanglement Structures (PSESs) (strict definition is explained later)

- undistinguishability
  - availability of verification of any maximally entangled states with small errors
  - an experimental verification with errors does not guarantee that the prepared state is exactly the same as the maximally entangled state
- self-duality ( $\Leftrightarrow \mathcal{K} = \mathcal{K}^*$ )
  - saturated situation of pre-duality ( $\Leftrightarrow \mathcal{K} \supset \mathcal{K}^*$ )
  - pre-duality  $\Leftarrow$  availability of projective measurements
- Assumption : Hereinafter, we assume  $\dim \mathcal{H}_A = \dim \mathcal{H}_B = d$

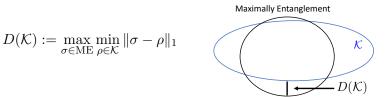
# Verification of maximally entangled states and $\epsilon$ -undistinguishability



- In order to prepare a maximally entangled state  $\sigma$ 
  - we prepare a state  $\rho$   $(\rho')$
  - $\blacktriangleright$  verify that  $\rho$   $(\rho')$  is equivalent to  $\sigma$  with error (probability) less than  $\epsilon$
- This task is successful if there exists such a state  $\rho, \rho'$  in a model.
  - $\blacktriangleright$  If  ${\cal K}$  has a state sufficiently near any maximally entangled state,  ${\cal K}$  success this task
  - ▶ e.g.  $\mathcal{K} \supset \mathcal{T}_+(\mathcal{H}_A \otimes \mathcal{H}_B)$  success this task

# Verification of maximally entangled states and $\epsilon$ -undistinguishability

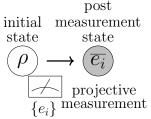
Because the error probability of verification (hypothesis testing) is estimated by trace norm  $\|\cdot\|_1$ , we introduce the following value for  $\mathcal{K}$ 



- if  $\mathcal K$  contains all maximally entangled states, then  $D(\mathcal K)=0$
- *e*-undistinguishability
  - $D(\mathcal{K}) \leq \epsilon$
  - ► if *K* is *ϵ*-undistinguishable, we cannot deny the possibility of *K* by any verification of maximally entangled states with larger errors than *ϵ*

## Relation between projectivity and pre-duality

- Projectivity
  - ► For any (pure) effect e, there exists projective measurement {e<sub>i</sub>} such that an element e<sub>i0</sub> is equal to e, and the post-measurement state is given as e<sub>i0</sub> := e<sub>i0</sub> / Tr e<sub>i0</sub>.
- $\mathcal{K}^*$  is generated by (pure) effects
- there exists a correspondence  $e_i \mapsto \overline{e_{i_0}}$
- because  $\overline{e_{i_0}}$  is a state,  $\overline{e_{i_0}} \in \mathcal{K}$
- $\rightarrow\,$  Projectivity derives pre-duality, i.e.,  $\mathcal{K}\supset\mathcal{K}^*$



- With preserving projectivity, state space is restricted and effect space is extended, then we obtain self-dual cone
  - (actually, this is not trivial and we show this statement later)

# The definition of $\epsilon\text{-}\mathsf{PSESs}$

- - ▶ *ϵ*-undistinguishability
  - self-duality

## Definition (*e*-Pseudo Standard Entanglement Structure)

If an ES  ${\cal K}$  satisfies  $\epsilon$ -undistinguishability and self-duality, we say that  ${\cal K}$  is an  $\epsilon$ -PSES.

- SES is an example of  $\epsilon\text{-PSES}$  for  $\epsilon\geq 0$
- No other simple example is not found
  - if  $\mathcal{K} \supseteq SES$ ,  $\mathcal{K}$  is not self-dual
  - if  $\mathcal{K} \subsetneq \text{SES}$ ,  $\mathcal{K}$  is not  $\epsilon$ -undistinguishable for small  $\epsilon$ .
- Q. Is there any other example of  $\epsilon$ -PSES for small  $\epsilon$  ?  $\rightarrow$  Yes!

#### Main Theorem

Given  $\epsilon > 0$ , there exists exactly different infinite models of  $\epsilon$ -PSES.

# Self-dual modification

#### Main Theorem 1

Given any pre-dual cone  $\mathcal{K}$ , there exists a self-dual cone  $\tilde{\mathcal{K}}$  such that  $\mathcal{K} \supset \tilde{\mathcal{K}} \supset \mathcal{K}^*$ .

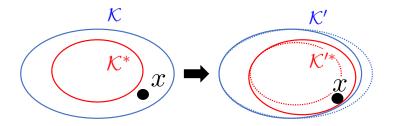
- Due to this theorem, self-duality is a natural consequence of projectivity.
- For proofs, we applying Zorn's Lemma for a certain ordered set (→The proof is not constructive)
- Different pre-dual cone cannot always be modified to different self-dual cone (This problem is solved by following theorem)

#### Main Theorem 2

Given an exact hierarchy of pre-dual cones  $\mathcal{K}_1 \supseteq \cdots \supseteq \mathcal{K}_n$ , there exist exactly different self-dual cones  $\tilde{\mathcal{K}}_i$  such that  $\mathcal{K}_i \supset \tilde{\mathcal{K}}_i \supset \mathcal{K}_i^*$ .

• pre-dual cones  $\mathcal{K}_1 \supsetneq \cdots \supsetneq \mathcal{K}_n \Rightarrow$  exactly different self-dual cones  $\tilde{\mathcal{K}_i}$ 

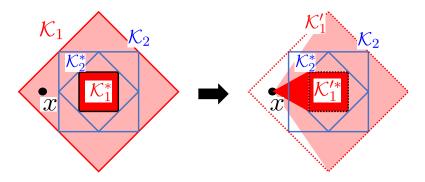
## Idea of Proof of Theorem 1



 $\begin{array}{c} \mathsf{Remark} \\ \mathcal{K} \supset \mathcal{K}' \Leftrightarrow \mathcal{K}^* \subset \mathcal{K}'^* \end{array}$ 

- if  $\mathcal{K} \subsetneq \mathcal{K}^*$ , there exists  $x \in \mathcal{K} \setminus \mathcal{K}^*$
- $\mathcal{K}'^* := \mathcal{K}^* + x$ ,  $\mathcal{K}' = (\mathcal{K}'^*)^*$
- then, x belongs to both  $\mathcal{K}'$  and  $\mathcal{K}'^*$
- $\rightarrow \ \mathcal{K} \subsetneq \mathcal{K}' \subset \mathcal{K}'^* \subsetneq \mathcal{K}^*$ 
  - repeat such step infinitely many times (by Zron's Lemma)

## Idea of Proof of Theorem 2



- $\mathcal{K}_1 \supseteq \mathcal{K}_2 \supseteq \mathcal{K}_2^* \supseteq \mathcal{K}_1^* \to \text{there exists } x \in \mathcal{K}_1 \setminus \mathcal{K}_2$
- $\mathcal{K}_1^{\prime*} := \mathcal{K}_1^* + x \rightarrow \mathcal{K}_1^{\prime*} \not\subset \mathcal{K}_2$   $\rightarrow \mathcal{K}_1^{\prime*} \subset \tilde{\mathcal{K}}_1^{\prime} \not\subset \tilde{\mathcal{K}}_2 \ ( \because \tilde{\mathcal{K}}_2 \subset \mathcal{K}_2)$  $\rightarrow \tilde{\mathcal{K}}_1^{\prime} \neq \tilde{\mathcal{K}}_2$

## The existence of PSESs

- In order to apply main theorem 2, we construct an exact hierarchy of pre-dual cones
- in this paper, we construct pre-dual cone  $\mathcal{K}_r$  for a parameter r>0
- $\mathcal{K}_r \supseteq \mathcal{K}_{r'}$  for  $r' < r \le r_0 \quad \rightarrow \quad {\{\mathcal{K}_r\}}$  is an exact hierarchy
- ightarrow Main Theorem 2 implies that  $ilde{\mathcal{K}_r}$  with each r is exactly different
  - $D(\tilde{\mathcal{K}}_r) \leq D(\mathcal{K}_r^*) \leq \epsilon$  for sufficiently small r.
- $\rightarrow \tilde{\mathcal{K}}_r$  is a  $\epsilon$ -PSES.

#### Main Theorem 3

Given  $\epsilon > 0$ , there exists exactly different infinite models of  $\epsilon$ -PSES.

## Other results

- an  $\epsilon$ -PSES has extraordinary performance for perfect discrimination
  - an  $\epsilon$ -PSES has non-orthogonal perfectly distinguishable states
  - ► In the SES, orthogonale ⇔ perfectly distinguishable

#### Main Theorem 4

For any  $\epsilon > 0$ , there is an  $\epsilon$ -PSES that contains a measurement discriminating two non-orthogonal states perfectly.

- Group symmetric condition characterize the SES uniquely
  - G-symmetric cone  $\mathcal{K}$ :  $g(x) \in \mathcal{K}$  for any  $x \in \mathcal{K}$  and any  $g \in G$ .

$$\operatorname{GU}(A;B) := \{g \in \operatorname{GL}(\mathcal{T}(\mathcal{H}_A \otimes \mathcal{H}_B)) \mid g(\cdot) := U^{\dagger}(\cdot)U,$$

U is a unitary matrix on  $\mathcal{H}_A \otimes \mathcal{H}_B$ .

#### Main theorem 5

For ES  $\mathcal{K}$ ,  $\mathcal{K}$  is  $\mathrm{GU}(A; B)$ -symmetric iff  $\mathcal{K} = \mathrm{SES}$ 

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## Summary and open problems

- There are many possibilities of ESs different from the SES.
- There exists infinite examples of  $\epsilon$ -PSESs
  - self-duality
  - *ϵ*-undistinguishability
- $\rightarrow$  There exists another possibility of ESs that cannot be denied by verification of maximally entangled state with errors
  - OPEN. Explicit construction of *e*-PSESs
  - Some  $\epsilon$ -PSES can discriminate non-orthogonal states
- $\rightarrow$  Verification of maximally entanglement state with errors cannot deny the possibility of an ES with such extraordinary performance for state discrimination

OPEN. Any  $\epsilon$ -PSES can discriminate non-orthogonal states except for the SES?

•  $\mathrm{GU}(A;B)$ -symmetry uniquely determines the SES OPEN. there exists  $\mathrm{LU}(A;B)$ -symmetric  $\epsilon$ -PSES except for the SES ?