

Computational self-testing for entangled magic states

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Beyondiid10 (2022)

Classical verifier



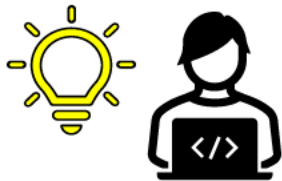
Quantum device (prover)



Goal: To know inner workings of black-box by classical verifier



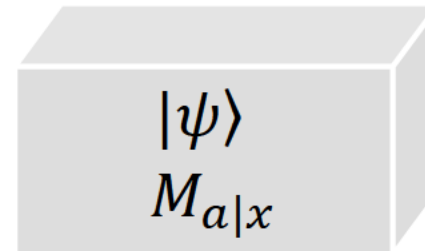
Classical verifier



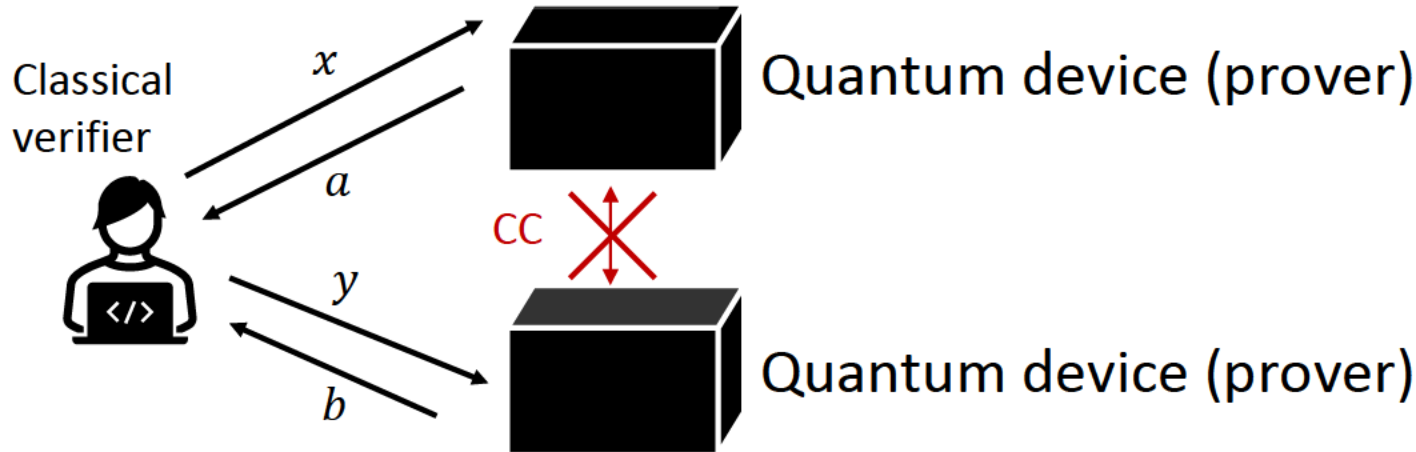
$x = 010$

$a = 110$

Quantum device (prover)



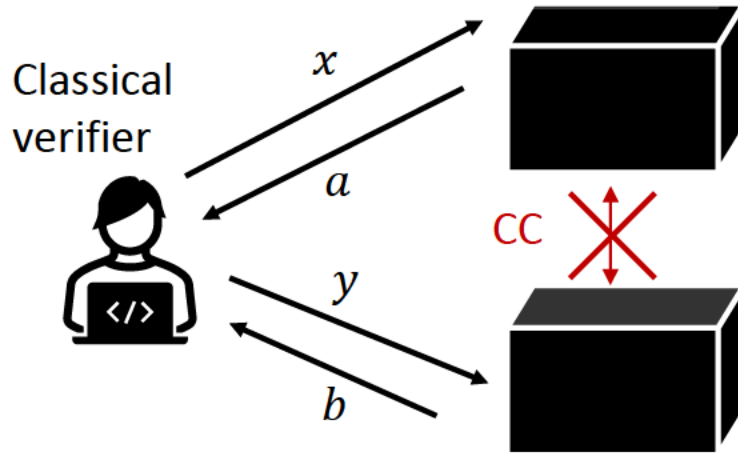
Standard self-testing



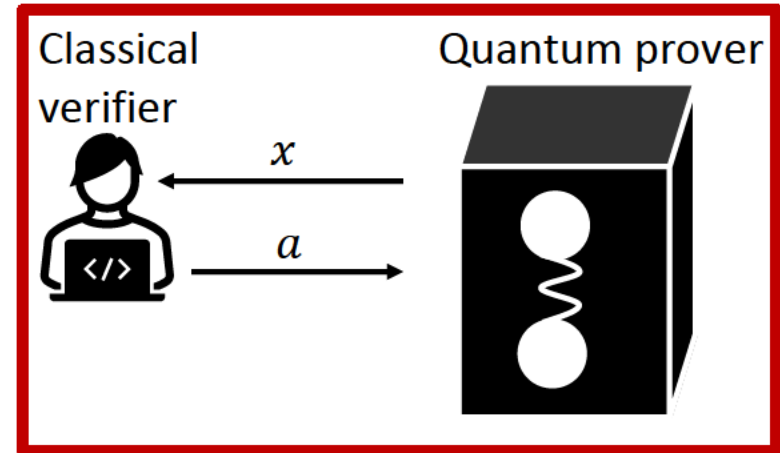
- Spatially-separated **multiple**-provers
- Assumption: Provers cannot communicate during the test
- \therefore To certify entanglement, execute test based on Bell inequality
- If devices **can** communicate, any classical devices can pass such self-testing protocol
- **Multi-prover assumption** is crucial

Self-testing with a single device

Standard self-testing



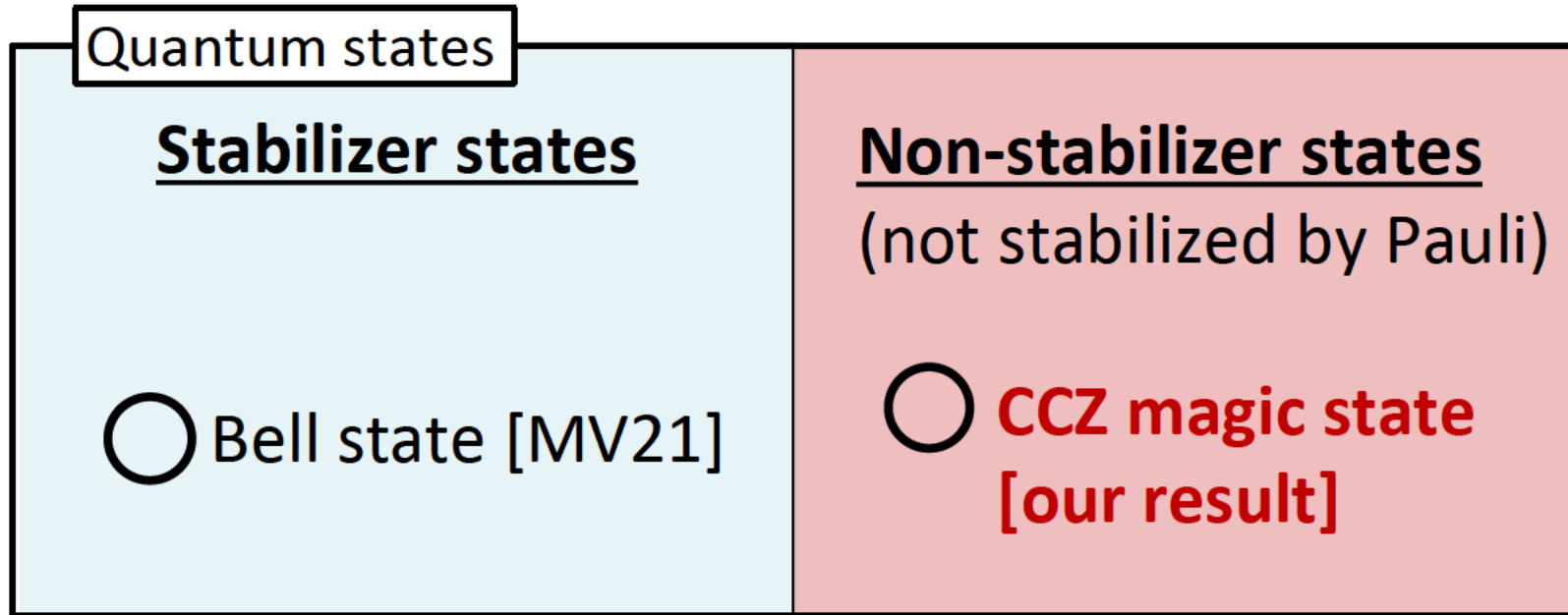
Metger, Vidick, Quantum 5,544 (2021)



- We can no longer rely on Bell inequality because any classical prover achieves the maximum violation
- [MV21] assumes a **quantum-poly-time prover** and cannot solve the Learning with errors (LWE) used in Post-Quantum cryptography
- **Self-test Bell state (stabilizer state)** relying on **stabilizer test**

Self-testing with a single device

- It should be straightforward to extend [MV21] to other stabilizer states
- We show that we can self-test **CCZ magic state (non-stabilizer state)**



1. Cryptographic primitive underlying the protocol
2. Our self-testing protocol for CCZ magic state

Trapdoor claw-free function families

Brakerski, et al, FOCS (2018), Mahadev, FOCS (2018)

From LWE assumption, function families F and G can be constructed

Function family $F = \{f_{k,0}, f_{k,1}\}_k$

1. 2 to 1:

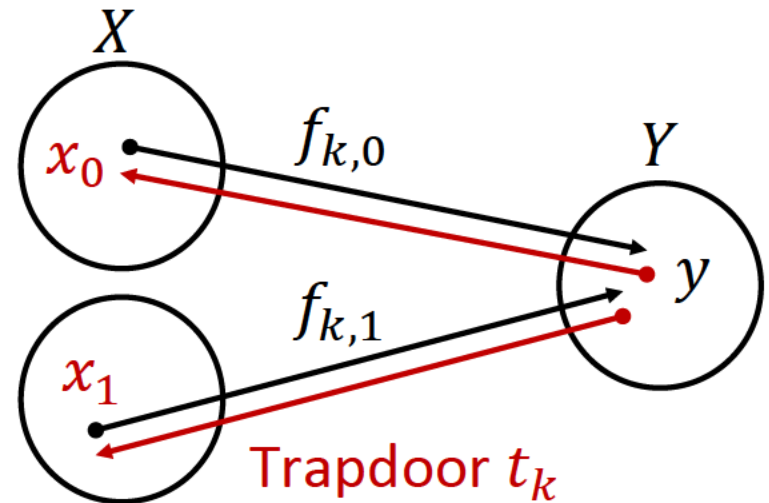
$$f: X \rightarrow Y$$

\exists preimages (x_0, x_1) of y s.t.

$$y = f_{k,0}(x_0) = f_{k,1}(x_1)$$

2. Claw-free property:

Quantum poly-time prover cannot calculate preimages (x_0, x_1) from y



$$x_b = \text{INV}_F(t_k, y, b)$$

$$b \in \{0,1\}$$

Trapdoor claw-free function families

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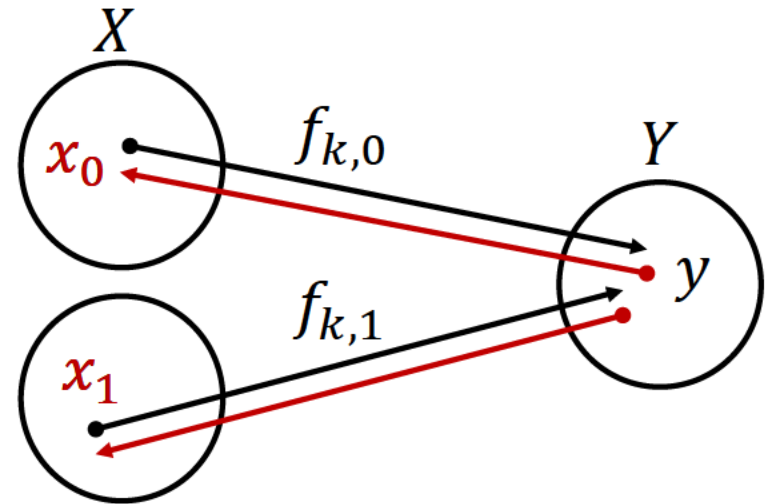
3. Adaptive hardcore bit property:

QPT prover is hard to output

$(y, b, x_b, r, d \neq 0)$

from k s.t.

$$y = f_{k,b}(x_b), \quad r = d \cdot (x_0 \oplus x_1) \in \{0,1\}$$



Hard to compute **one preimage** and **any-one bit information** about the other preimage

Trapdoor claw-free function families

Brakerski, et al, FOCS (2018), Mahadev, FOCS (2018)

From LWE assumption, function families F and G can be constructed

Function family $G = \{f_{k,0}, f_{k,1}\}_k$

1. 1 to 1:

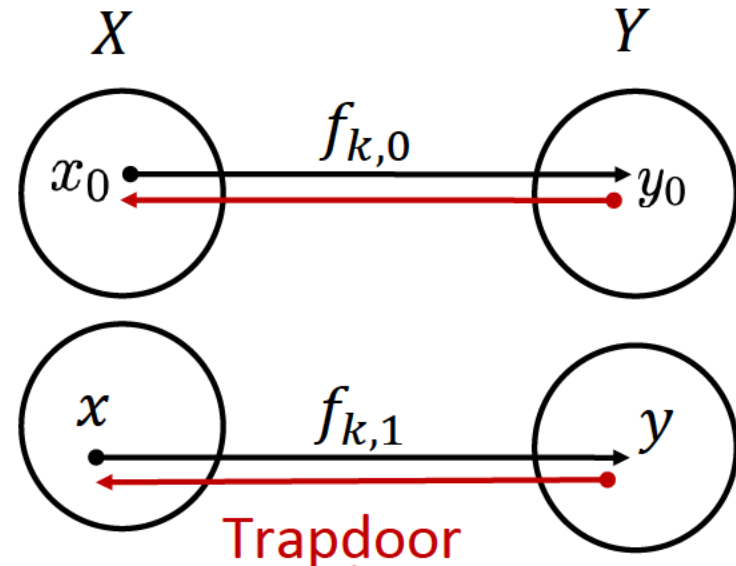
$$f: X \rightarrow Y$$

\exists preimage x of y s.t.

$$y = f_{k,b}(x)$$

2. Injective invariance:

Quantum poly-time prover is hard to distinguish between F and G from public key k



$$(1, x) = INV_G(t_k, y)$$

Using F and G , we construct our self-testing protocol

Choose $\theta_1 \theta_2 \theta_3 \in \{000, 001, 010, 100, 111\}$
 $\theta = 1 \rightarrow$ generate (k, t_k) of F
 $\theta = 0 \rightarrow$ generate (k, t_k) of G

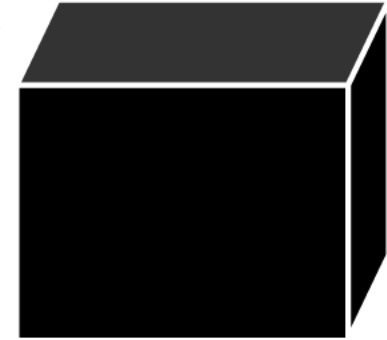
From k , I cannot find θ



Classical verifier

Public keys k_1, k_2, k_3

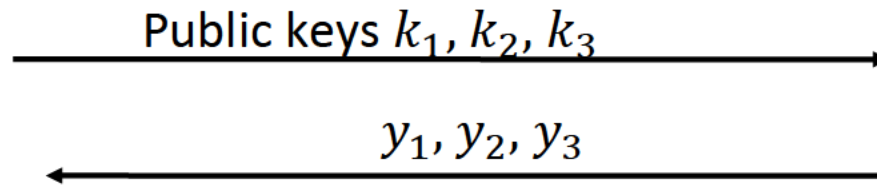
Quantum prover



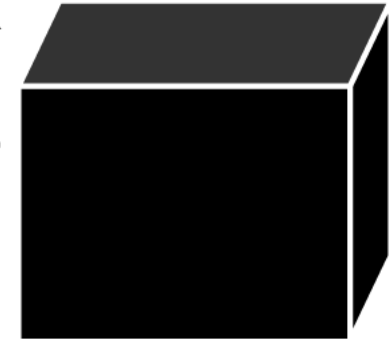
$$\sum_{b=0}^1 \sum_{x \in X} |b\rangle |x\rangle |f_{k,b}(x)\rangle$$

Choose $\theta_1 \theta_2 \theta_3 \in \{000, 001, 010, 100, 111\}$
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Quantum prover



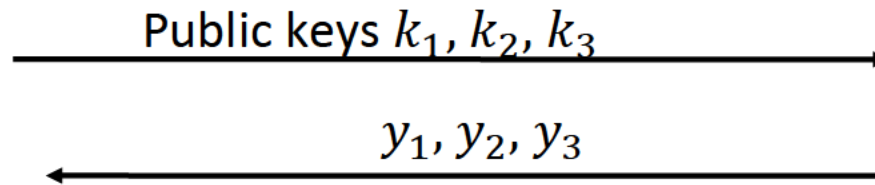
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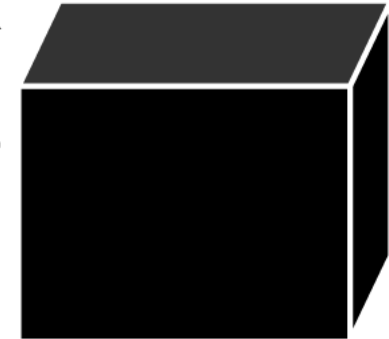
Self-testing of CCZ magic state

Choose $\theta_1 \theta_2 \theta_3 \in \{000, 001, 010, 100, 111\}$
 $\theta = 1 \rightarrow$ generate (k, t_k) of F
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From k , I cannot find θ



Quantum prover



$$\theta = 0 \rightarrow |b\rangle|x_b\rangle$$
$$\theta = 1 \rightarrow |0\rangle|x_0\rangle + |1\rangle|x_1\rangle$$

Trapdoor claw-free function families

Brakerski, et al, FOCS (2018), Mahadev, FOCS (2018)

From LWE assumption, function families F and G can be constructed

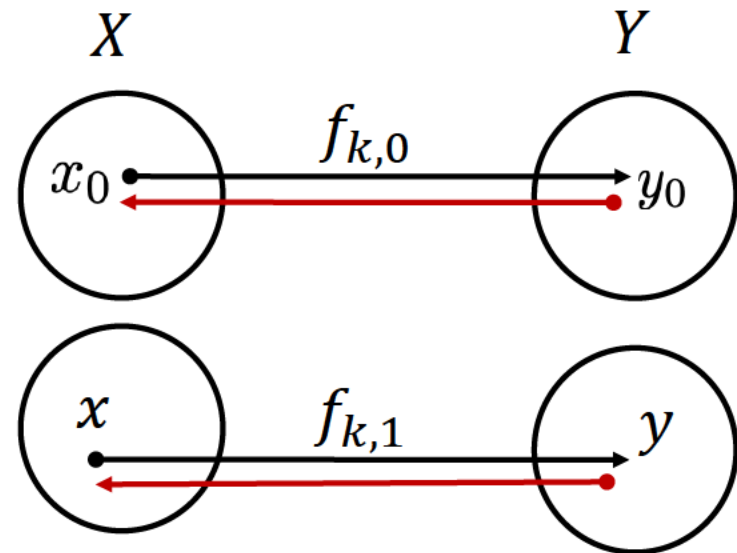
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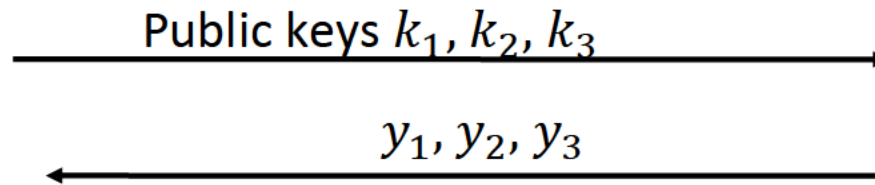
\exists preimage x of y s.t.

$$y = f_{k,b}(x)$$

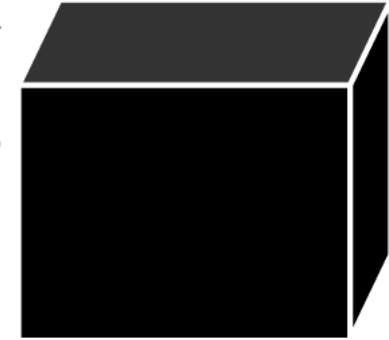


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Quantum prover



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$$\theta = 1 \rightarrow |0\rangle|x_0\rangle + |1\rangle|x_1\rangle$$

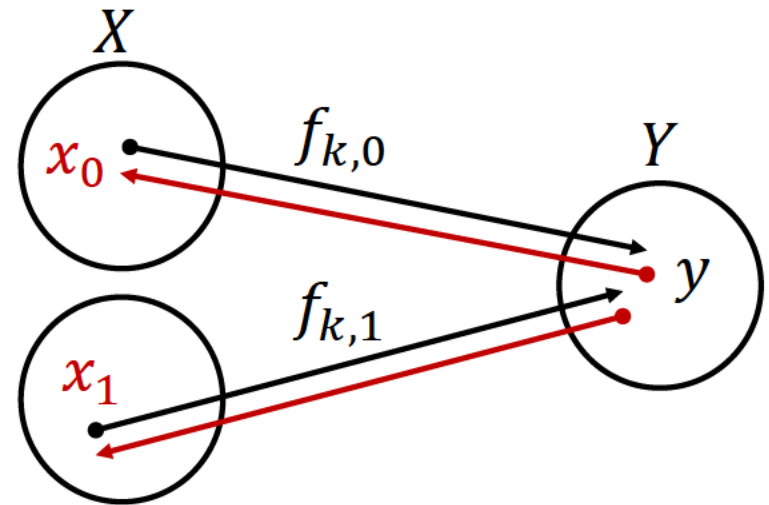
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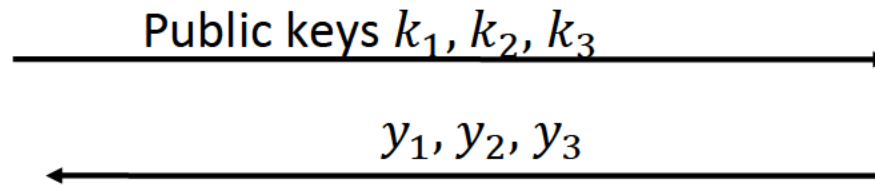
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Quantum prover



$$\theta = 0 \rightarrow |b\rangle|x_b\rangle$$
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From k , I cannot find θ



Classical verifier

Public keys k_1, k_2, k_3

y_1, y_2, y_3

Z measurement

(b, x_b)

Quantum prover



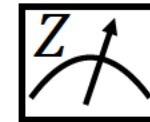
Check of Z-measurement:

Check if (b, x_b) is correct
 $(y = f_{k,b}(x_b))$ holds or not

Failure prob. P_Z

$$\theta = 0 \rightarrow |b\rangle|x_b\rangle$$

$$\theta = 1 \rightarrow |0\rangle|x_0\rangle + |1\rangle|x_1\rangle$$

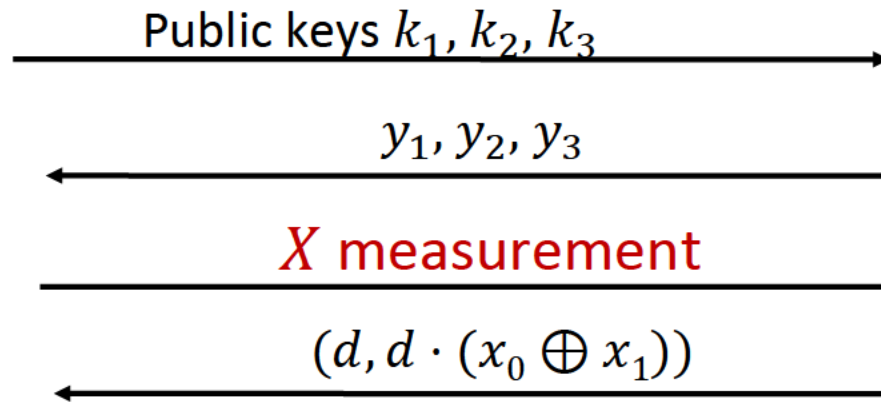


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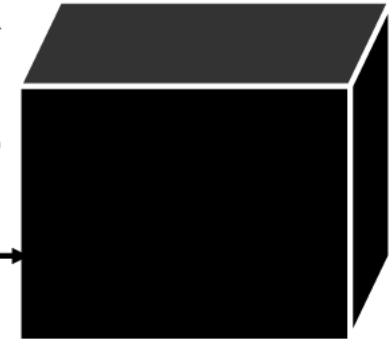
From k , I cannot find θ



Classical verifier



Quantum prover



Check of X-measurement :

Using trapdoor, check if outcomes are equal to $(d, d \cdot (x_0 \oplus x_1))$

Failure prob. P_X

$$\theta = 1 \rightarrow |0\rangle|x_0\rangle + |1\rangle|x_1\rangle$$



$$|0\rangle + (-1)^{d \cdot (x_0 \oplus x_1)} |1\rangle$$

$$m = d \cdot (x_0 \oplus x_1)$$

Trapdoor claw-free function families

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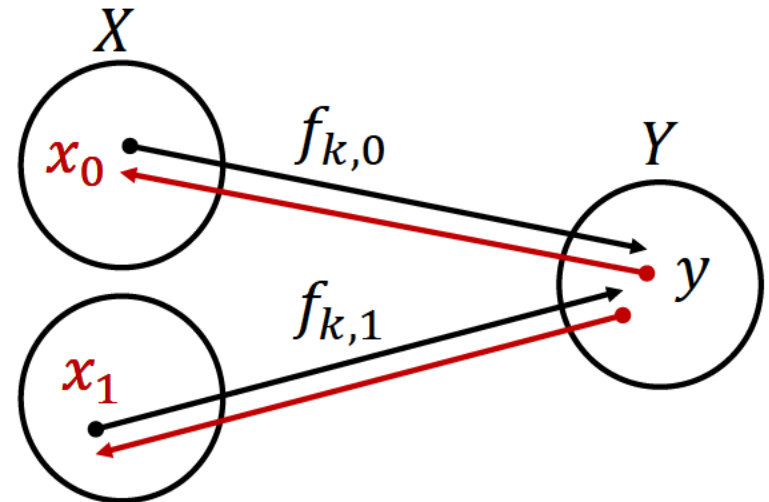
3. Adaptive hardcore bit property:

QPT prover is hard to output

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from k s.t.

$$y = f_{k,b}(x_b), r = d \cdot (x_0 \oplus x_1) \in \{0,1\}$$



Honest quantum prover satisfies only one of the two relations, but cannot by **any** quantum prover due to the hardness of LWE

Choose $\theta_1 \theta_2 \theta_3 = 111$



Public keys k_1, k_2, k_3

$y_1, y_2, y_3, d_1, d_2, d_3$

Generalized stabilizer test

(v_1, v_2, v_3)

Quantum prover



Using trapdoor, check if (v_1, v_2, v_3) is correct
Failure prob. P_{STAB}

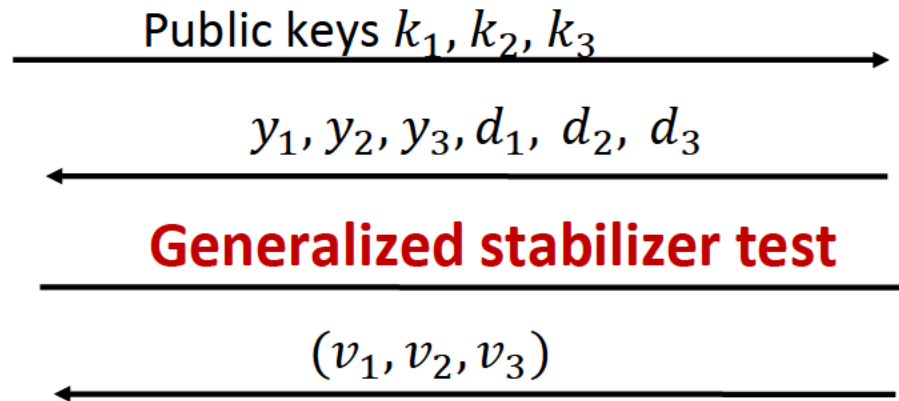
$$|CCZ\rangle = CCZ \left[\bigotimes_{i=1}^3 (|0\rangle + (-1)^{m_i} |1\rangle) \right]$$

$$m = d \cdot (x_0 \oplus x_1)$$

Choose $\theta_1 \theta_2 \theta_3 = 111$



Classical verifier



Quantum prover



Any pure state $|\psi\rangle = U|0^n\rangle$ is stabilized by Hermitian operators:

$$\text{Generalized stabilizers } \{g_i\}_{i=1}^n = \{UZ_iU^\dagger\}_{i=1}^n$$

$$g_1|CCZ\rangle = (X \otimes CZ)|CCZ\rangle = |CCZ\rangle$$

Generalized stabilizer test: Check measurement outcomes of $\{g_i\}_{i=1}^n$

Is a 2-qubit gate needed to measure the outcome of $X \otimes CZ$?

Choose $\theta_1 \theta_2 \theta_3 = 111$



Classical verifier

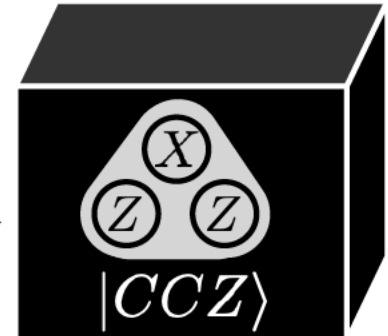
Public keys k_1, k_2, k_3

$y_1, y_2, y_3, d_1, d_2, d_3$

Generalized stabilizer test

(v_1, v_2, v_3)

Quantum prover



Takeuchi, Morimae, PRX (2018)

$$(X \otimes CZ) |CCZ\rangle = \left(X \otimes \sum_i |i\rangle\langle i| \otimes Z^i \right) |CCZ\rangle = |CCZ\rangle$$

\downarrow \downarrow \downarrow
 v_1 v_2 v_3

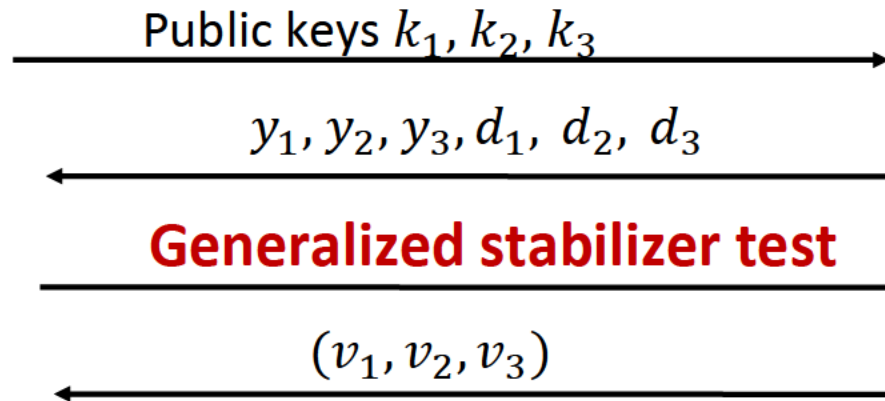
+1 eigenspace of $X \otimes CZ \Leftrightarrow v_1 \oplus v_2 v_3 = m$

Only the verifier knows

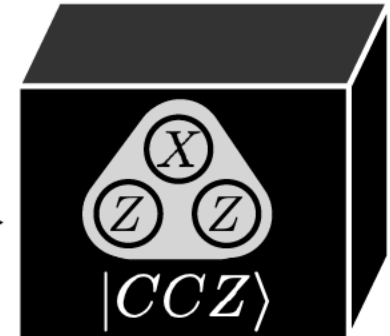
Choose $\theta_1 \theta_2 \theta_3 = 111$



Classical verifier



Quantum prover



Takeuchi, Morimae, PRX (2018)

$$(X \otimes CZ) |CCZ\rangle = \left(X \otimes \sum_i |i\rangle\langle i| \otimes Z^i \right) |CCZ\rangle = |CCZ\rangle$$

\downarrow \downarrow \downarrow
 v_1 v_2 v_3

- Using **Z and X measurement checks**, verifier knows how accurately prover performs **generalized stabilizer measurements**
- By checking $\{v_i\}_i$, verifier finds the distance of ρ and $|CCZ\rangle$

Theorem

- Prover's failure probabilities of Z, X basis measurements and generalized stabilizer measurements: $(P_Z, P_X, P_{\text{STAB}})$
- Prover's Hilbert space H

There exists an isometry $V: H \rightarrow C^8 \otimes H, \zeta_H$ and constant $r > 0$ such that for general prover's state ρ when $\theta = 111$, the following holds:

- Self-testing of state

$$\|V\rho V^\dagger - |CCZ\rangle\langle CCZ| \otimes \zeta_H\|^2 \leq O(P_Z^r + P_X^r + P_{\text{STAB}}^r)$$

- Self-testing of state and measurement

$$|\text{tr}[\rho P_{\vec{q}}^{(\vec{v})}] - |\langle \vec{v}_{\vec{q}} | CCZ \rangle|^2| \leq O(P_Z^r + P_X^r + P_{\text{STAB}}^r)$$

Projector of outputting $\vec{v} = v_1 v_2 v_3$
given request $\vec{q} = q_1 q_2 q_3 \in \{Z, X\}^3$

$|000\rangle$ if $\vec{q} = ZZZ, \vec{v} = 000$
 $|+++ \rangle$ if $\vec{q} = XXX, \vec{v} = 000$

Our contribution

Using Trapdoor claw-free function families, we reveal that the **non-stabilizer state $|CCZ\rangle$ can also be self-tested with a single-quantum device**

Future work

1. Self-test states where X , Y and Z appear in generalized stabilizers
 - If only X and Z appear, its single-device self-testing seems possible
 - It seems difficult as we have Trapdoor claw-free function families composed of only **2** function families F and G
2. Self-test in non-iid scenario
 - We assume prover's behavior is the same for each repetition
 - Is it possible by using de Finetti's theorem or by one-shot self-testing?